

# Why Haskell Does Not Matter

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# About Me

- Started programming in Haskell as a CS undergraduate at Utrecht University (1998).
- MSc Software Technology (2005).
- PhD studies in the field of type-based program analysis (2005–2010).
- Software Development Engineer at Vector Fabrics (since March 1, 2010).



**Universiteit Utrecht**



**Vector**Fabrics

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# Vector Fabrics

- High-tech start-up, founded in February 2007.
- Unique technology for designing and implementing multicore/multithreaded systems.
- 15 employees.
- Funded with venture capital from the US and the Netherlands.
  
- vfAnalyst:
  - First released product (March 1st, 2010).
  - On-line source-code analysis.
- Future products focus on program transformation and mapping to hardware.

- Customers
  - log in at `www.vectorfabrics.com`;
  - upload C source files and an accompanying test set through a web interface.
- vfAnalyst
  - runs and analyses the code;
  - visualises properties of interest;
  - pinpoints performance bottlenecks;
  - uncovers opportunities for parallelisation;
  - generates reports.

The screenshot displays the vfAnalyst web interface, divided into two main sections. The left section shows a coverage report table, and the right section shows a call graph visualization.

**Coverage Report Table:**

Name	Coverage	Load	Actual	Target
exit	0	0.00	0.00	0.00
fopen	100	0.00	0.00	0.00
read_rgb	100	6.86	0.00	0.00
Loop_1	100	6.86	0.00	0.00
perfor	0	0.00	0.00	0.00
exit	0	0.00	0.00	0.00
fclose	100	0.00	0.00	0.00
fclose	100	0.00	0.00	0.00
Loop_21	100	93.13	0.00	0.00
printf	100	0.00	0.00	0.00
<b>rgb2yu</b>	<b>100</b>	<b>27.56</b>	<b>0.00</b>	<b>0.00</b>
Loop	100	27.56	0.00	0.00
write_y	100	34.31	0.00	0.00
read_rg	100	31.26	0.00	0.00

**Properties for rgb2yuv:**

- Function: [rgb2yuv](#)
- Compute load: 27.6 %
- Line coverage: 100.0 %
- Source location: [rgb2yuv.c: 9-29](#)
- Call location: [main.c 34](#)

**Call Graph:**

The call graph shows the execution flow starting from `f_main` through `f_read_rgb` to `Loop_21`, which then calls `rgb2yuv`. `rgb2yuv` calls `write_yuv`, which in turn calls `Loop_24`. The `Loop_24` node is expanded to show its internal structure, including a `write_yuv` call and several other sub-loops represented by green arrows.

- Combination of static and dynamic analysis.
- C compiler that targets a propriety assembly format.
- Small virtual machine executes assembly and obtains precise information about
  - how different parts of the program communicate through different memory locations;
  - execution times of different functions;
  - code coverage.



# Practices

- Short release cycles split up in short (2-to-4-week) development iterations.
- Daily stand-up meetings.
- Collective code ownership.
- Firm coding standards.
- Distributed SCM.
- Pair programming.
- Frequent code reviews.
- Continuous integration.
- Fully automatic unit tests.
- Fully automatic acceptance/regression tests.

We don't use Haskell.

- Choice for Caml was made already before the company was officially founded.
- Established research contact between Utrecht University and Philips Research sparked an interest in functional programming amongst the founders.
- Of the languages considered (Caml and Haskell), Haskell was recognised as the more mature and the elegant.
- Still, **eager evaluation** tipped the balance to Caml.
  - In some experiments performance of Haskell programs was very hard to predict.
  - In these experiments Caml performed much more consistently.
- Concern for a small company: how do we recruit Caml programmers?

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# Many Different Programming Languages

- GUI: Flex/ActionScript.
- Database management: Caml.
- Static analyses: Caml.
- Dynamic analyses: C.
- C compiler: Caml, C++ (LLVM).
- Virtual machine: C.
- Content-management system: PHP (Joomla).
  
- A fair dose of Python and shell scripts to glue everything together.
  
- Each of these components communicates data, often to different machines.

- Caml has a good foreign-function interface to C... but not much else.
  - To enable communication between components written in Caml and the Flex GUI, Caml functions need to be exposed through a C wrapper.
  - From the C wrapper, a tool called SWIG (Simplified Wrapper and Interface Generator) these C wrappers can be interfaced with in Python.
  - The Python wrappers, finally, can be immediately called from Flex GUI.

# Data Marshalling

- With all this indirection, we really need to keep the amount of effort spent in marshalling data to a minimum.
- Caml functions that are exposed to the GUI only take simple types, such as integers and strings, as input.
- They produce either simple types or XML.
- The produced XML is either processed by the server-side Python scripts or forwarded to the client-side Flex GUI (resulting in quite a lot of bandwidth usage).
  
- Maintaining, extending, or modifying the communication between components is **difficult** and **hard to debug**.

# Database Access

- Similiar issues arise in communication with our database.
- Both the Caml and Python components interact with a single MySql database.
- However, the bindings for these two languages make different assumptions about the behaviour of the database.



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# A Embedded-system Designer's Perspective

- Unfamiliar to functional programming before joining Vector Fabrics.
- Reasonable easy to learn Caml.
- Biggest hurdle: adapt to the “functional mindset”:
  - use recursion rather than for-loops;
  - use higher-order functions;
  - exploit polymorphism.
- Biggest grievance: very poor type-error messages.
- Main advantages: productivity, easier-to-maintain code.
- Main drawback: the price paid for high-level abstractions.

# A Haskell Programmer's Perspective

- We adapted to Caml very quickly.
- We really miss having type signatures in the main source files.
- We really miss freely (re)ordering function definitions and type declarations.
- We make some use of the camlp4 preprocessor to extend Caml's syntax.
- We make good use of the module system to organise software both in the small and the large:
  - compared to Haskell type classes, the explicitness of modules is sometimes a gain, sometimes a loss;
  - the structuring of modules sometimes get quite complex.
- We deliberately limit the use of mutable references.
- We think of **eagerness** as a mixed blessing. [▶ Some code](#)

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# Summary

- Trying to build a successful company that uses functional programming for its core development tasks is lots of fun.
- We need a couple of success stories in order to convince the next high-tech start-up to adapt Haskell.

# Acknowledgements



- Alexey Rodriguez Yakushev
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# Evaluation Order

```
let f x y = () (* TODO:... *)
let prerr_usage = prerr_string "Usage: foo arg1 arg2\n"
let main =
  if Array.length Sys.argv ≤ 2 then (* check #args *)
  begin prerr_usage; exit 1 end else (* too few: print error *)
  let x = Sys.argv.(1) in (* enough: do some work *)
  let y = Sys.argv.(2) in
  f x y
```

```
$ ./foo 2 3 5 7 11
Usage: foo arg1 arg2
$
```

# Searching a List

## A Haskeller's Take

```
(* hd :  $\alpha$  list  $\rightarrow$   $\alpha$  *)
```

```
let hd = function x :: _  $\rightarrow$  x
```

```
(* filter : ( $\alpha \rightarrow$  bool)  $\rightarrow$   $\alpha$  list  $\rightarrow$   $\alpha$  list *)
```

```
let rec filter p = function
```

```
  | []  $\rightarrow$  []
```

```
  | x :: xs  $\rightarrow$  if p x then x :: filter p xs else filter p xs
```

```
(* find : ( $\alpha \rightarrow$  bool)  $\rightarrow$   $\alpha$  list  $\rightarrow$   $\alpha$  *)
```

```
let find p xs = hd (filter p xs)
```



# Searching a List

## Using Exceptions

```
exception Not_found
```

```
(* iter : ( $\alpha \rightarrow$  unit)  $\rightarrow$   $\alpha$  list  $\rightarrow$  unit *)
```

```
let rec iter f = function
```

```
| []  $\rightarrow$  ()
```

```
| x :: xs  $\rightarrow$  f x; iter f xs
```

```
exception Found of my_type
```

```
(* find : (my_type  $\rightarrow$  bool)  $\rightarrow$  my_type list  $\rightarrow$  my_type *)
```

```
let find p xs =
```

```
  try
```

```
    iter (fun x  $\rightarrow$  if p x then raise (Found x) else ()) xs;
```

```
    raise Not_found
```

```
  with Found x  $\rightarrow$  x
```

# Searching a List

## Using a Dedicated Function

```
(* find : ( $\alpha$  → bool) →  $\alpha$  list →  $\alpha$  *)
```

```
let rec find p = function
```

```
| [] → raise Not_found
```

```
| x :: xs → if p x then x else find p xs
```

# What About Searching a Tree?

## Using Exceptions

```
type  $\alpha$  tree = Leaf of  $\alpha$  | Node of  $\alpha$  tree  $\times$   $\alpha$  tree
```

```
(* find : ( $\alpha$   $\rightarrow$  bool)  $\rightarrow$   $\alpha$  tree  $\rightarrow$   $\alpha$  *)
```

```
let rec find p = function
```

```
  Leaf x  $\rightarrow$  if p x then x else raise Not_found
```

```
  Node (l, r)  $\rightarrow$  try find p l with Not_found  $\rightarrow$  find p r
```

# What About Searching a Tree?

## Using Optional Values

```
type  $\alpha$  tree = Leaf of  $\alpha$  | Node of  $\alpha$  tree  $\times$   $\alpha$  tree  
type  $\alpha$  option = None | Some of  $\alpha$ 
```

```
(* find : ( $\alpha$   $\rightarrow$  bool)  $\rightarrow$   $\alpha$  tree  $\rightarrow$   $\alpha$  *)  
let find p t =  
  let rec find_aux = function  
    | Leaf x  $\rightarrow$  if p x then Some x else None  
    | Node (l,r)  $\rightarrow$  match find_aux p l with  
      | None  $\rightarrow$  find_aux p r  
      | x_opt  $\rightarrow$  x_opt  
  in  
  match find_aux p t with  
    | None  $\rightarrow$  raise Not_found  
    | Some x  $\rightarrow$  x
```

# What About Searching a Tree?

## Using Continuations

```
type  $\alpha$  tree = Leaf of  $\alpha$  | Node of  $\alpha$  tree  $\times$   $\alpha$  tree
```

```
(* find : ( $\alpha$   $\rightarrow$  bool)  $\rightarrow$   $\alpha$  tree  $\rightarrow$   $\alpha$  *)
```

```
let find p =
```

```
  let rec find_aux k = function
```

```
    | Leaf x  $\rightarrow$  if p x then x else k ()
```

```
    | Node (l,r)  $\rightarrow$  find_aux (fun _  $\rightarrow$  find_aux k r) l
```

```
in
```

```
  find_aux (fun _  $\rightarrow$  raise Not_found)
```

◀ Return