

Random Testing of Purely Functional Abstract Datatypes

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Abstract datatypes

- Defined only by their operations
- Independent from a concrete implementation
- Implementations can change without affecting client codes
- Client codes can easily switch between implementations

Algebraic specification

- Fitting framework for the definition of abstract datatypes
- In particular in the context of *purely* functional languages
- Enables equational reasoning

Equational reasoning

- Substituting “equals for equals”
- Deriving a whole class of theorems from only a handful of axioms
- Implementors only need to make sure that the axioms hold

Property-based random testing

- Map axioms to testable properties
- Obtain an arbitrary large set of executable test cases
- Excellent fit for purely functional languages
- QuickCheck, Gast, ...

Algebraic specification

Equational reasoning

Property-based random testing

Test cases

Theorem₁
⋮
Theorem_m

Implementation₁
⋮
Implementation_n



Algebraic specification

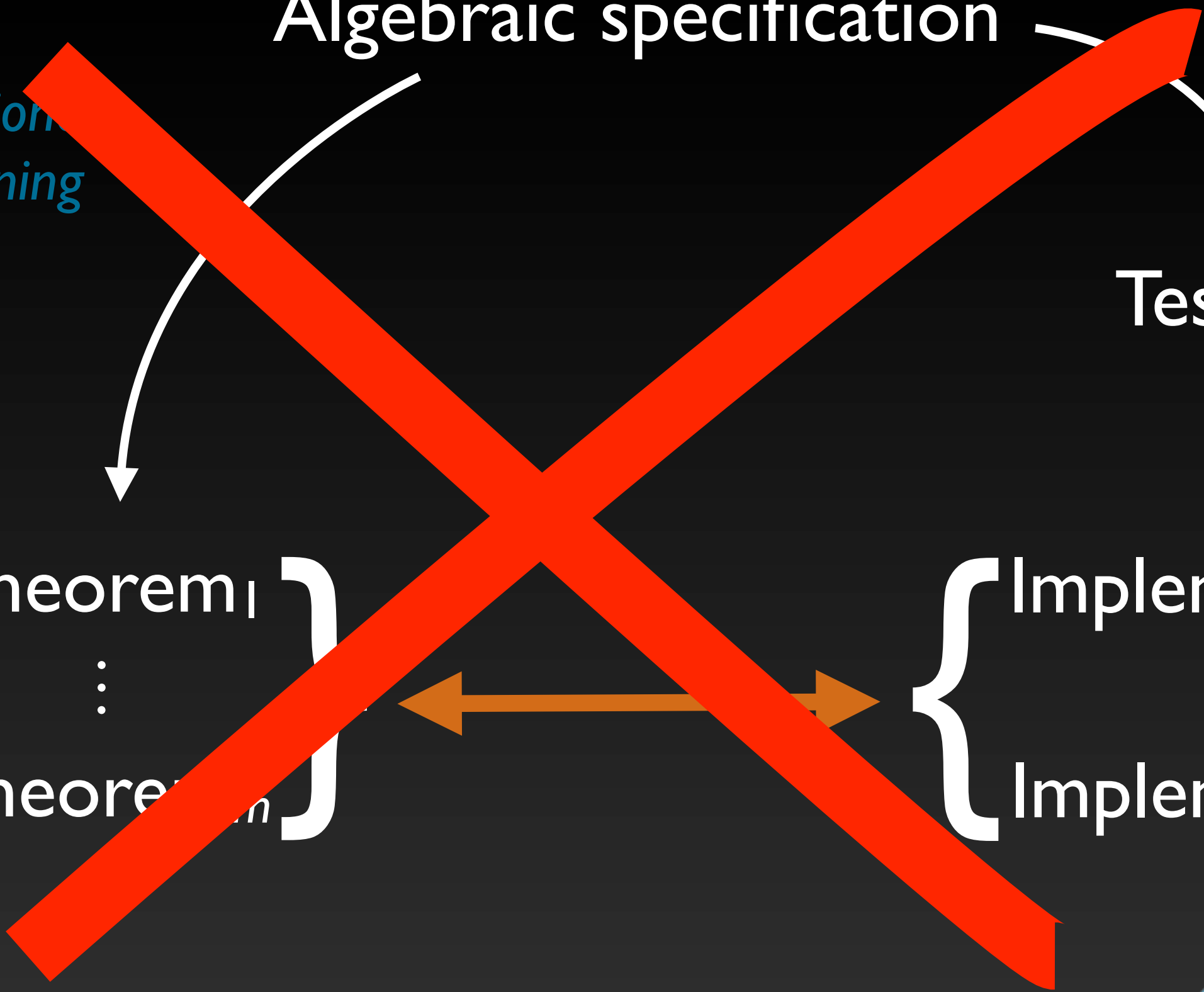
Equation reasoning

Property-based random testing

Test cases

Theorem₁
⋮
Theorem_n

Implementation₁
⋮
Implementation_n



Koopman et al. (IFL 2011)

“Without detailed study of the internals of the [implementation of an] ADT it is undecidable if a set of logical properties is sufficient [to assert its correctness with respect to the specification].”

Outline

- A failing example
- What went wrong?
- A solution
- Conclusion

A failing example

FIFO queues: *signature*

sort:

Queue

operations:

```
empty      :: Queue
enqueue    :: Int → Queue → Queue
isEmpty    :: Queue → Bool
front      :: Queue → Int
dequeue    :: Queue → Queue
```

FIFO queues: *axioms*

Q1: `isEmpty empty` = `True`

Q2: `isEmpty (enqueue x q)` = `False`

Q3: `front (enqueue x empty)` = `x`

Q4: `front (enqueue x q)` = `front q` (if `isEmpty q = False`)

Q5: `dequeue (enqueue x empty)` = `empty`

Q6: `dequeue (enqueue x q)` = `enqueue x (dequeue q)` (if `isEmpty q = False`)

A theorem about queues

```
isEmpty (dequeue (enqueue x empty)) = True
```

Proof:

```
isEmpty (dequeue (enqueue x empty))  
= { Q5 }  
isEmpty empty  
= { Q1 }  
True
```

Another theorem about queues

```
front (dequeue
  (enqueue x (enqueue y (enqueue z empty)))) = y
```

Proof:

```
front (dequeue (enqueue x (enqueue y (enqueue z empty))))
= { Q6, Q2 ; Q6, Q2 }
front (enqueue x (enqueue y (dequeue (enqueue z empty))))
= { Q5 ; Q4, Q2 }
front (enqueue y empty)
= { Q3 }
y
```

QuickCheck by example

```
> let p1 = property (\xs ys →  
    reverse (xs ++ ys) == reverse ys ++ reverse xs)
```

```
> quickCheck p1  
+++ OK, passed 100 tests
```

```
> let p2 = property (\x → reverse [x] == [])
```

```
> quickCheck p2  
*** Failed! Falsifiable (after 1 test):  
0
```

Testable properties for queues

```
q1 = property (isEmpty empty)
q2 = property ( $\lambda x q \rightarrow \neg(\text{isEmpty } (\text{enqueue } x q))$ )
q3 = property ( $\lambda x \rightarrow \text{front } (\text{enqueue } x \text{ empty}) == x$ )
q4 = property ( $\lambda x q \rightarrow \neg(\text{isEmpty } q) ==> \text{front } (\text{enqueue } x q) == \text{front } q$ )
q5 = property ( $\lambda x \rightarrow \text{dequeue } (\text{enqueue } x \text{ empty}) == \text{empty}$ )
q6 = property ( $\lambda x q \rightarrow \neg(\text{isEmpty } q) ==>$ 
    dequeue (enqueue x q) == enqueue x (dequeue q))
```


Testable properties for queues

```
q1 = property (isEmpty empty)
q2 = property ( $\lambda x q \rightarrow \neg(\text{isEmpty } (\text{enqueue } x q))$ )
q3 = property ( $\lambda x \rightarrow \text{front } (\text{enqueue } x \text{ empty}) == x$ )
q4 = property ( $\lambda x q \rightarrow \neg(\text{isEmpty } q) ==> \text{front } (\text{enqueue } x q) == \text{front } q$ )
q5 = property ( $\lambda x \rightarrow \text{dequeue } (\text{enqueue } x \text{ empty}) == \text{empty}$ )
q6 = property ( $\lambda x q \rightarrow \neg(\text{isEmpty } q) ==>$ 
    dequeue (enqueue x q) == enqueue x (dequeue q))
```

Batched queues

```
data Queue = BQ [Int] [Int] deriving Show
```

```
bq [] r = BQ (reverse r) []
```

```
bq f r = BQ f r
```

```
empty = bq [] []
```

```
enqueue x (BQ f r) = bq f (x : r)
```

```
isEmpty (BQ f r) = null f
```

```
front (BQ f r) = last f
```

```
dequeue (BQ f r) = bq (tail f) r
```

Batched queues

```
data Queue = BQ [Int] [Int] deriving Show
```

```
bq [] r = BQ (reverse r) []
```

```
bq f r = BQ f r
```

```
empty = bq [] []
```

```
enqueue x (BQ f r) = bq f (x : r)
```

```
isEmpty (BQ f r) = null
```

```
front (BQ f r) = last f -- incorrect!
```

```
dequeue (BQ f r) = bq (tail f) r
```

Equality for batched queues

```
instance Eq Queue where
```

```
  q1 == q2 = toList q1 == toList q2
```

```
toList (BQ f r) = f ++ reverse r
```

Testing batched queues

```
> mapM_ quickCheck [q1,q2,q3,q4,q5,q6]
+++ OK, passed 100 tests.
+++ OK, passed 100 tests.
+++ OK, passed 100 tests.
+++ OK, passed 100 tests.
+++ OK, passed 100 tests.
+++ OK, passed 100 tests.
```

What went wrong?

One more property...

```
> let q7 = property ( $\lambda x y z \rightarrow$   
    front (dequeue (enqueue x (enqueue y (enqueue z empty)))) == y)
```

```
> quickCheck q7
```

```
*** Failed! Falsifiable (after 2 tests):
```

```
0
```

```
1
```

```
0
```

But... q7 represents one of our theorems!

Did we break equational reasoning?

Yes

Did we break equational reasoning?

- A stable basis for equational reasoning requires that operations are *invariant* under equality:

$$x = y \Rightarrow h(x) = h(y)$$

- Invariance is what justifies “substituting equals for equals”

Did we break equational reasoning?

```
> let qA = BQ [2,3] [5]
```

```
> let qB = BQ [2] [5,3]
```

```
> qA == qB
```

```
True
```

```
> front qA
```

```
3
```

```
> front qB
```

```
2
```

A solution

Key idea

Systematically extend the set of testable properties
with properties for operation invariance

A first attempt

```
> let qq = property ( $\lambda q q' \rightarrow$   
     $q == q' \wedge \neg(\text{isEmpty } q) ==>$   
     $\text{front } q == \text{front } q'$ )
```

```
> quickCheck qq
```

```
*** Gave up! Passed only 1 test.
```

A type of equivalent values

```
data Equiv a = a :: a deriving Show
```

For example:

```
> let eq = BQ [2,3] [5] :: BQ [2] [5,3]
```

(More details in the paper)

Lifting out the equivalence check...

```
> let qq' = property ( $\lambda(q ::= q')$  →  
  ¬(isEmpty q) ==>  
  front q == front q')
```

Lifting out the equivalence check...

```
> let qq' = property ( $\lambda(q ::=: q') \rightarrow$   
   $\neg(\text{isEmpty } q) \implies$   
   $\text{front } q == \text{front } q')$ 
```

```
> quickCheck qq'
```

```
*** Failed! Falsifiable (after 4 tests):
```

```
BQ [-1,-2] [2] ::=: BQ [-1] [2,-2]
```


Testable invariance properties

```
qq1 = property (λx (q ==: q') → enqueue x q == enqueue x q' )
qq2 = property (λ  (q ==: q') → isEmpty q == isEmpty q' )
qq3 = property (λx (q ==: q') → ¬(isEmpty q) ==> front q == front q' )
qq4 = property (λx (q ==: q') → ¬(isEmpty q) ==> dequeue q == dequeue q' )
```

Testing against all properties

```
> mapM_ quickCheck ([q1,q2,q3,q4,q5,q6] ++ [qq1,qq2,qq3,qq4])
+++ OK, passed 100 tests.
+++ OK, passed 100 tests.
+++ OK, passed 100 tests.
+++ OK, passed 100 tests.
+++ OK, passed 100 tests.
+++ OK, passed 100 tests.
+++ OK, passed 100 tests.
+++ OK, passed 100 tests.
+++ OK, passed 100 tests.
*** Failed! Falsifiable (after 5 tests):
BQ [-1,0] [1] ==: BQ [-1] [1,0]
+++ OK, passed 100 tests.
```

Conclusion

Summary

- A framework for tests and proofs for purely functional ADTs
- An extension for dealing with implementations that are not UR
- Key idea: derive testable properties for operation invariance

Future work

- Assess and quantify impact on real-world applications
- Automatic derivation of testable properties from specifications
- EDSL for algebraic specifications of ADTs
- Testable properties for Equiv -generators